

The lattice of linear Mal'cev conditions

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Identities of height at most 1 are usually called **linear**.

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Not examples

group terms, lattice terms, semilattice term.

Definition (Barto, Pinsker)

An algebra \mathbf{A} is said to be a **reflection** of \mathbf{B} defined by mappings $h_1: B \rightarrow A$ and $h_2: A \rightarrow B$, if for every basic operation f we have

$$f_{\mathbf{A}}(a_1, \dots, a_n) = af_{\mathbf{B}}(b(a_1), \dots, b(a_n)).$$

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The class of all retractions (reflections, resp.) of algebras from \mathcal{K} is denoted $\mathbf{R}\mathcal{K}$ ($\mathbf{R}_{\text{ret}}\mathcal{K}$).

Theorem (Barto, Pinsker, O)

A class of algebras is definable by linear identities (identities of height 1, resp.) if and only if it is closed under \mathbf{R}_{ret} and \mathbf{P} (\mathbf{R} and \mathbf{P} , resp.).

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A variety of groups is congruence regular, uniform and also singular.

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Congruence singularity ...

Characterization by linear identities

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(Except congruence singularity.)

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Meet of Mal'cev conditions

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Meet of two Mal'cev conditions is the strongest Mal'cev condition that which is weaker than both of the original ones.

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Observation

A clone satisfies this Mal'cev condition if and only if it is a product of a clone with Mal'cev operation and a clone with Jónsson terms.

A **product of clones** \mathcal{A} and \mathcal{B} is the clone \mathcal{C} with $C = A \times B$, and $\mathcal{C}^{[n]} = \{f \times g : f \in \mathcal{A}^{[n]}, g \in \mathcal{B}^{[n]}\}$.

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Proof.

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Thank you for your attention!